

2024

Information theory and Coding

Full Marks: 100

Time: Three hours

*The figures in the margin indicate full marks for the questions.**Answer any five questions.*

1.	a)	Define a discrete memory less source (DMS). What is meant by the entropy of a DMS? Show that the average information content of the DMS is given by: $H(S) = \sum_{i=0}^{M-1} p_i \times \log_2 \left(\frac{1}{p_i} \right)$; where ' p_i ' is the probability of the occurrence of the RV ' x_i ' of a random process S.	2+2+6
	b)	Consider the transformation of a RV ' Y ' by assuming the RV ' X ' has a pdf $f_X(x) = e^{-x} \times u(x)$; where ' $u(x)$ ' is a unit step Heaviside function. If the relation between the RVs ' Y ' and ' X ' is $Y = X^3$; show that the pdf of ' Y ' is given by $f_Y(y) = \frac{e^{-\sqrt[3]{y}}}{3} \times \sqrt[3]{\frac{1}{y^2}} \times u(y)$.	10
2.	a)	b) Prove that the pdf of the sum of two statistically independent RV ($W = X + Y$) is the convolution of their individual RVs.	10
	b)	Show that the mutual information between the transmitted symbol ' x_i ' and the received symbol ' y_j ' is given by $I(x_i y_j) = \log_2 \left[\frac{p(x_i y_j)}{p(x_i)} \right]$. Hence show that the average mutual information of the channel is given by $I(X;Y) = H(X) - H(X Y)$; where the symbols have their usual meaning.	4+6
3.	a)	Show that the entropy of an extended DMS (S^n) is given by $H(S^n) = n \times H(S)$; where ' n ' is the number of symbols in each block of the DMS with an alphabet size ' M '. A DMS has an alphabet $\{S_0, S_1\}$ with marginal probabilities $p(S_0) = p_0 = \frac{1}{4}$ and $p(S_1) = p_1 = \frac{3}{4}$. Find the entropies of the source ' S ' and the extended source ' S^3 '.	5+5
	b)	A source produces three symbols 'A, B and C' with the following marginal and conditional probabilities:	3+7

		<table border="1"> <tr> <td>i</td> <td>$p(i)$</td> </tr> <tr> <td>A</td> <td>1/4</td> </tr> <tr> <td>B</td> <td>1/4</td> </tr> <tr> <td>C</td> <td>1/2</td> </tr> </table>	i	$p(i)$	A	1/4	B	1/4	C	1/2	<table border="1"> <tr> <td>$p(j i)$</td> <td>A</td> <td>B</td> <td>C</td> </tr> <tr> <td>A</td> <td>1/8</td> <td>1/4</td> <td>5/8</td> </tr> <tr> <td>B</td> <td>1/2</td> <td>1/8</td> <td>3/8</td> </tr> <tr> <td>C</td> <td>3/8</td> <td>5/8</td> <td>0</td> </tr> </table>	$p(j i)$	A	B	C	A	1/8	1/4	5/8	B	1/2	1/8	3/8	C	3/8	5/8	0	
i	$p(i)$																											
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C	3/8	5/8	0																									
		<p>(i) Assume that there is no inter-symbol interference; calculate the entropy of the source.</p> <p>(ii) Calculate the conditional entropy $H(Y X)$</p>																										
4.	a)	<p>Prove that when the random variable 'x' corresponds to a Gaussian distribution having mean 'a' and probability density function</p> $f_X(x) = \frac{1}{\sigma\sqrt{2\pi}} \times \exp\left[-\frac{(x-a)^2}{2\sigma^2}\right]$ <p>the differential entropy will be given by</p> $h(X) = \frac{1}{2} \log_2(2\pi e\sigma^2)$ <p>where 'σ^2' is the variance of the Gaussian process 'X'.</p>		10																								
	b)	<p>Show that the differential entropy '$h(X)$' of a continuous random variable is a maximum when the probability density function of the random variable is Gaussian, and the maximum value of the differential entropy is given by</p> $h(X) = \frac{1}{2} \log_2(2\pi e\sigma^2)$ <p>where the symbols have their usual meaning.</p>		10																								
5.	a)	<p>Discuss Shannon's information capacity theorem. Hence, prove that the channel capacity of a band-limited, power-limited Gaussian channel with input power 'S' is given by $C = B \times \log_2\left(1 + \frac{S}{N}\right)$ bits/sec, and the symbols have their usual meaning.</p>		4+10																								
	b)	<p>What is an ideal information system? Calculate Shannon's limit for an ideal system.</p>		6																								
6.	Write short notes on any two from the following:		10+10																									
	a)	Average mutual information is always positive ($I(X;Y) \geq 0$)																										
	b)	Binary symmetric channel (BSC) and transitional probability.																										
	c)	Channel capacity of a BSC: $C_s _{BSC} = [1 - H(p)]$ bits/sym.																										
	d)	Optimum distribution of input power.																										

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